

algorithms for solving semidefinite programs

- interior point methods
- spectral bundle methods
- bundle method
- projection methods

semidefinite programs: primal and dual

$$(SDP) \begin{cases} \min & \langle C, X \rangle \\ \text{s.t.} & \mathcal{A}(X) = b \\ & X \succeq 0 \end{cases}$$

 $\min_{X \succeq 0} \max_{y \in \mathbb{R}^m} \langle C, X \rangle + \langle b - \mathcal{A}(X), y \rangle \geq \max_{y \in \mathbb{R}^m} \min_{X \succeq 0} \langle b, y \rangle + \langle X, C - \mathcal{A}^\top(y) \rangle$

$$(\mathbf{DSDP}) \begin{cases} \max & b^{\top}y \\ \text{s.t.} & \mathcal{A}^{\top}(y) + Z = C \\ & y \in \mathbb{R}^m, Z \succeq 0 \end{cases}$$

SDP for max-cut and ϑ -number

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- →
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- max-cut: sizes of interest *n* around 100, strengthened relaxation leads to more than 10 000 constraints.
- ∂-number: sizes of interest n ≥ 500, results in 100 000 constraints.
- \rightarrow need other algorithmic machinery than interior point methods.

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assume m and n are large \longrightarrow avoid Cholesky factorization, matrix multiplication,...
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idea: get rid of $Z \succeq 0$ by using eigenvalue arguments.

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constant trace property holds for many SDP derived from combinatorial optimization problems.

Reformulate dual as follows:

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 \longrightarrow used to compute dual variable λ explicitely

$$\lambda_{min}(Z^*) = 0 \iff \lambda_{max}(-Z^*) = 0$$

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$$\min\{b^{ op}y + \lambda_{\max}(\mathcal{C} - \mathcal{A}^{ op}(y)) \colon y \in \mathbb{R}^m\}$$

 \rightarrow non-smooth unconstrained convex problem in y.

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 \longrightarrow can be done by iterative methods even for very large (sparse) matrices.

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two ingredients:

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with new $k \times k$ matrix variable V.

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This is a strictly convex function, if u > 0 is fixed.

substituting the definition of \hat{f} ...

$$\min_{y\in\mathbb{R}^m}\hat{f}(y)+\frac{u}{2}\|y-\hat{y}\|^2=$$

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quadratic SDP in the $k \times k$ matrix variable V, since $W = PVP^{\top}$. k is user defined and can be small, in particular, it is independent of n.

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Having point y_{new} , evaluation $f(y_{new})$ (sparse eigenvalue computation) produces also an eigenvector v to λ_{max} .

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 \rightarrow eigenvector v is added as new column to P, and P is purged by removing unnecessary other columns.

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software SBmethod a C++ implementation of the spectral bundle method of [Helmberg and Rendl 00; Helmberg and Kiwiel 99] no longer supported, but now there is the ConicBundle callable library instead; available at https:

//www-user.tu-chemnitz.de/~helmberg/ConicBundle/

example: consider again the basic max-cut relaxation $\max\{\langle L,X\rangle\colon {\rm diag}(X)=e,X\succeq 0\}$

Now 20 000 $\leq n \leq$ 50 000, sparse graphs.

п	upper-bnd	cut	time (secs)
20,000	143.3	131.3	330
20,000	261.9	244.8	536
20,000	598.1	571.1	1255
30,000	214.9	197.2	753
30,000	393.3	367.4	990
30,000	897.9	857.3	2330
40,000	286.9	262.7	1180
40,000	524.6	489.8	1650
50,000	358.9	328.5	1800

spectral bundle method summarized

- using eigenvalue optimization and classical methods from convex analysis
- general tool for solving SDP having matrices of large dimension
- convergence is slow, once close to optimum



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$$z^* = \max\{\langle C, X \rangle \colon \mathcal{A}(X) = a, \ \mathcal{B}(X) = b, \ X \succeq 0\}$$

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$$\mathcal{L}(X; y) = \langle C, X \rangle + y^{\top}(b - \mathcal{B}(X))$$

dual functional:

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basic assumption: we can evaluate f(y) easily, yielding also a maximizer X^* and $g^* = b - \mathcal{B}(X^*)$.

using
$$g^* = b - \mathcal{B}(X^*)$$
 if X^* is the optimizer for given \bar{y} :
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hence $g^* = b - \mathcal{B}(X^*)$ is subgradient.

since

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$$z^* = \min_{y \in \mathbb{R}^m} f(y) \le f(\tilde{y}) \quad \forall \tilde{y} \in \mathbb{R}^m$$

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any $\tilde{y} \in \mathbb{R}^m$ provides upper bound on z^* \longrightarrow try to find tight upper bound (i.e., approximate minimizer of f(y)) by using bundle methods.

two ingredients:

- work with a "bundle" of X_i's and maximize over conv{X₁,...,X_k} instead of over {A(X) = a, X ≥ 0}
- ▶ penalize displacement from current iterate, i.e., add penalty term $\frac{1}{2t} \|y \hat{y}\|^2$

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$$\begin{split} \min_{y} \hat{f}(y) + \frac{1}{2t} \|y - \hat{y}\|^2 \\ \hat{f}(y) &= \max_{X \in \text{conv}\{X_1, \dots, X_k\}} \mathcal{L}(X; y) \end{split}$$

iterative procedure:

solve approximately

$$\min_{\gamma\geq 0}\hat{f}(y) + \frac{1}{2t}||y - \hat{y}||^2$$

where

$$\hat{f}(y) = \max\{\mathcal{L}(X;y) \colon X \in \mathsf{conv}\{X_1,\ldots,X_k\}$$

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computational effort in each iteration:

- solve a convex quadratic program in k variables
- evaluate $f(\hat{y})$ to yield new \hat{X} and a subgradient \hat{g}

example: solving the max-cut relaxation strengthened by triangle-inequalities:

$$\max\{\langle L, X \rangle \colon \operatorname{diag}(X) = e, \mathcal{M}(X) \leq -e, X \succeq 0\}$$

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SDP with matrix dimension n and only n linear constraints.

bundle method

SDP relaxation for max-cut; triangle inequalities are dualized. 50 bundle iterations for n = 800, and 30 for n = 2000.

graph	п	initial gap (%)	final (%)	time (secs)
G6	800	22.29	18.15	43.11
G11	800	11.56	1.54	60.20
G14	800	4.51	2.84	59.68
G18	800	18.38	7.96	69.19
G22	2000	6.34	5.66	278.06
G27	2000	25.77	22.94	406.66
G39	2000	21.27	12.63	533.36

see [Fischer, Gruber, Rendl, Sotirov, 06]

software: ConicBundle C++ library of Ch. Helmberg, available at
https:
//www-user.tu-chemnitz.de/~helmberg/ConicBundle/

bundle method summarized

- in combination with interior point methods is a good tool to approximate SDPs with a huge number of constraints
- the number of function evaluations to reach good approximations is surprisingly small
- getting to the "real" optimum is hard



algorithms for solving semidefinite programs

- interior point methods
- spectral bundle method
- bundle methods
- projection methods

min f(x) such that $x \in \mathcal{X}$, h(x) = 0

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 $f : \mathbb{R}^n \mapsto \mathbb{R}, h : \mathbb{R}^n \mapsto \mathbb{R}^m$ sufficiently smooth functions, $\mathcal{X} \subseteq \mathbb{R}^n$ nonempty closed convex set of simple structure

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repeat until convergence (a) Keep y fixed: solve $\min_x \mathcal{L}_{\sigma}(x, y)$ to get x (b) update y: $y \leftarrow y + \sigma h(x)$ (c) update σ

Original version: Powell, Hestenes, 1969

(**DSDP**) min
$$b^{\top}y$$
 s.t. $\mathcal{A}^{\top}(y) - C = Z, Z \succeq 0$

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$$\mathcal{L}_{\sigma}(y, Z; X) = b^{ op} y + \langle X, Z + C - \mathcal{A}^{ op}(y)
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inner minimization

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- ...

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$$= b^\top y + \frac{1}{\sigma} \|Z - \mathcal{W}(y)\|^2 - \frac{1}{2\sigma} \|X\|^2$$

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$$\min_{y,Z\succeq 0} b^\top y + \frac{1}{\sigma} \|Z - \mathcal{W}(y)\|^2$$

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 $V \succeq 0, Z \succeq 0, VZ = 0.$

solve coordinatewise:

keep Z (and X) constant, y is given by the unconstrained minimization

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boundary point method

Initialization: k = 0, select $\sigma_k > 0$, $X_k \succeq 0, Z_k \succeq 0$

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See Malick, Povh, Rendl, W., 2007

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example: when computing the ϑ -number: $\mathcal{A}(\mathcal{A}^{\top}(.))$ is diagonal. Therefore, the main computational effort is the projection on the positive semidefinite cone.

example: computing the $\vartheta\text{-number}$ for random graphs from the Kim Toh collection

graph	п	E	time (secs)
theta82	400	23871	87
theta83	400	39861	70
theta102	500	37466	143
theta103	500	62515	110
theta104	500	87244	124
theta123	600	90019	205
theta162	800	127599	570

boundary point method summarized

- works "orthogonal" to interior point methods
- convergence behavior not well understood
- for matrices of moderate size, but can deal with a large number of constraints

Semidefinite Programming solvers at the NEOS site: http://www.neos-server.org/

- csdp
- dsdp
- penbmi
- pensdp
- sdpa
- sdplr
- sdpt3
- sedumi

moreover: sdplib [B. Borchers] at http://euler.nmt.edu/~brian/sdplib/ and the sdp website [Ch. Helmberg] at http://www-user.tu-chemnitz.de/~helmberg/semidef.html

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thank you for your attention!